

Rook-drawing for plane graphs

Claire Pennarun

David Auber, Nicolas Bonichon and Paul Dorbec

LaBRI, Bordeaux

Graph Drawing
September 25th, 2015

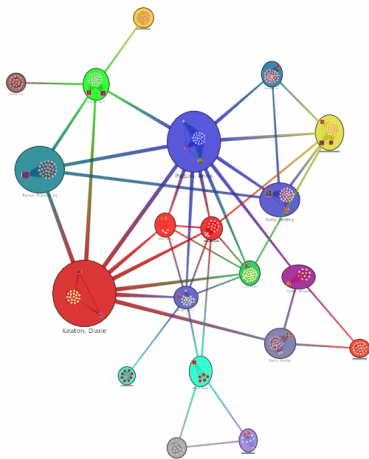
MAYBE YOU KNOW THEM?



DRAWING GRAPHS

We want to draw **large** graphs with hierarchical view: a vertex in the drawing = a group of vertices in the graph

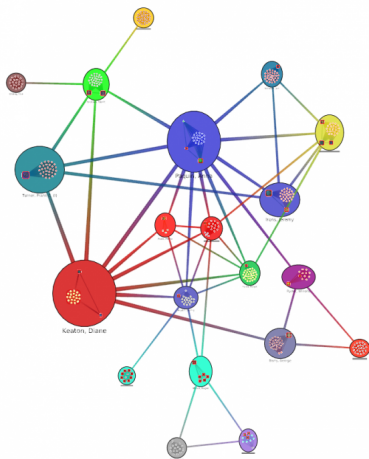
- preserve the relative positions of vertices → mental map
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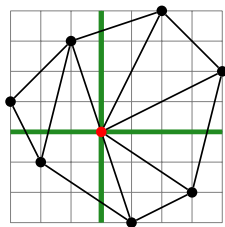
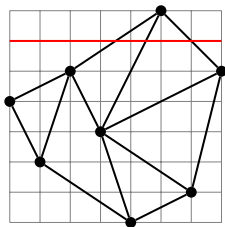
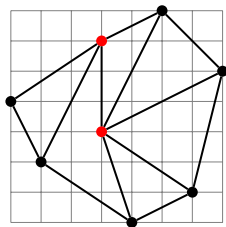


→ new type of drawing with constraints: rook-drawing

ROOK-DRAWING

A **rook-drawing** of a graph of n vertices:

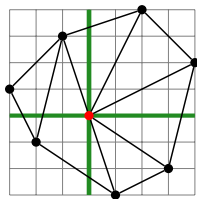
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- Regular grid (n columns and rows)
- Exactly one vertex per row and column



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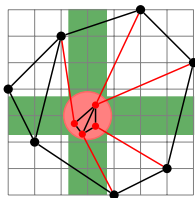
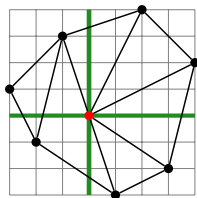
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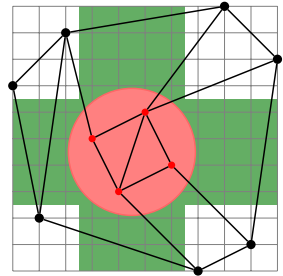
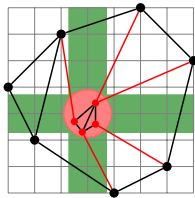
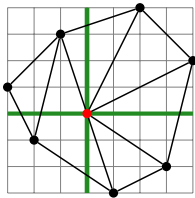
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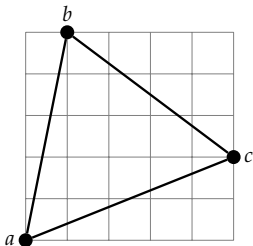
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What we already know:

- Straight-lines drawing ([Fáry, 1948] : every planar graph)
- Grid drawing ([de Fraysseix, 1988], [Schnyder, 1990] : every plane graph on an $(n - 2) \times (n - 2)$ grid)

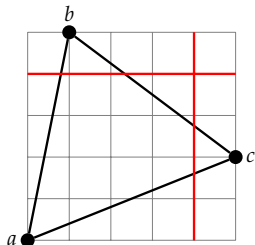
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Three exterior vertices a , b and c , n vertices (here $n = 6$).



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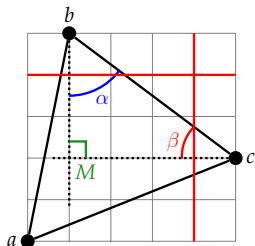
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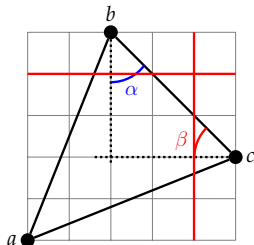
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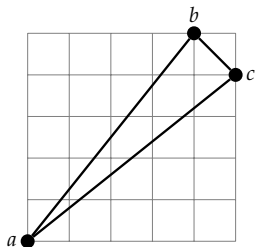
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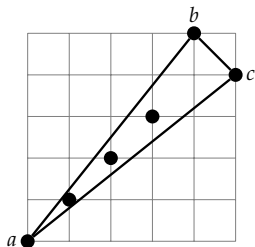
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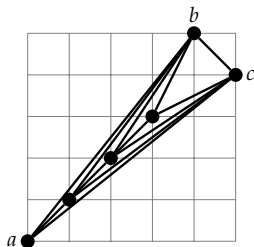
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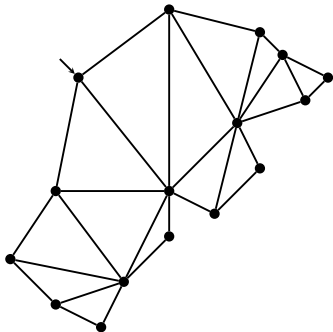
Graph with a degree 3 outer face with planar rook-drawing = subgraph of the tower graph

ROOK-DRAWING FOR OUTERPLANAR GRAPHS

A graph is **outerplanar** if it has a planar drawing such that all its vertices are on the outer face.

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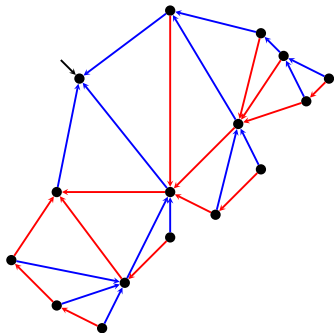


[Bonichon, Gavoille, Hanusse, 2005]

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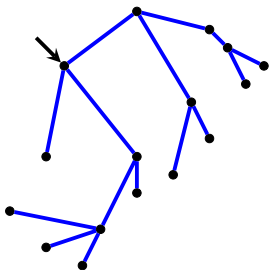


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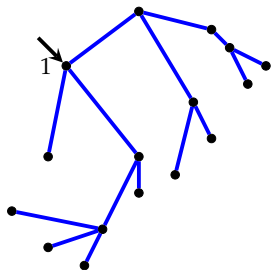
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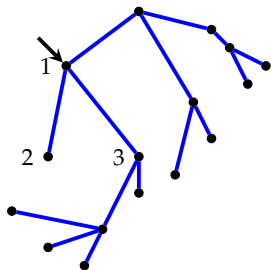
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- x : ccw pre-order depth-first
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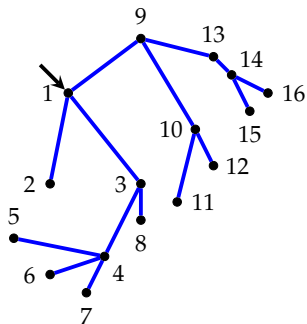
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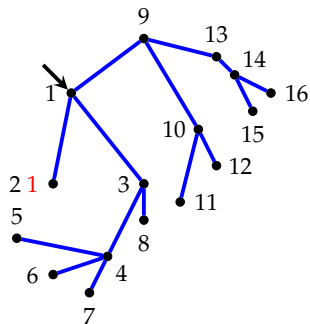
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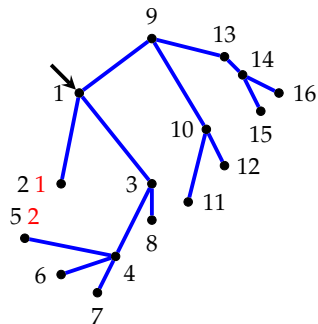
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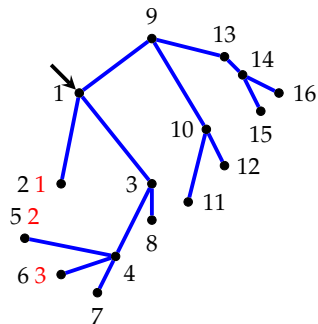
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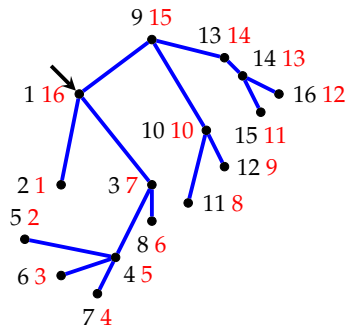
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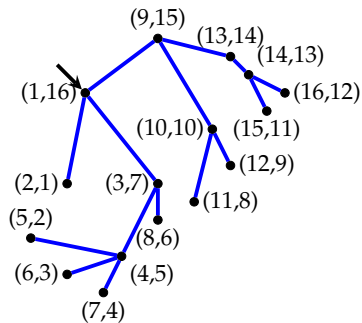
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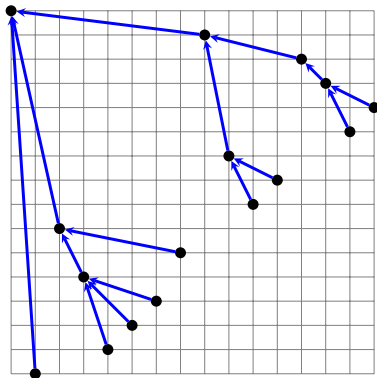
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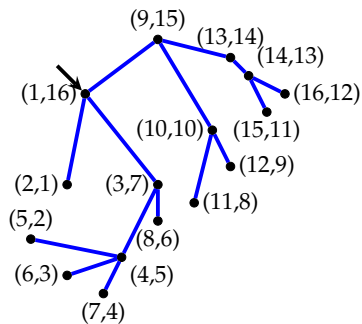


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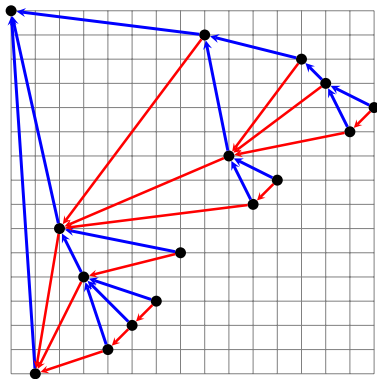


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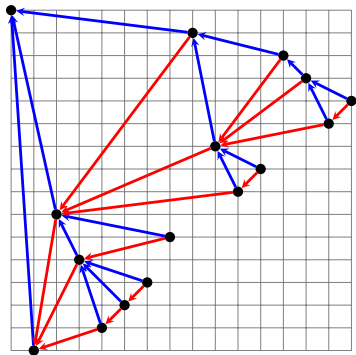


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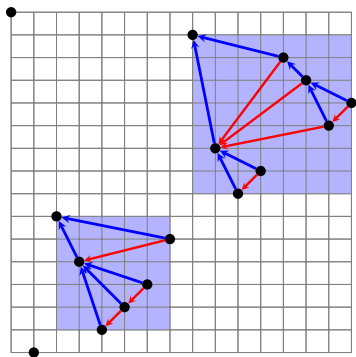
Hyp: $G[T_v]$ of depth k admits a planar rook-drawing in the grid $[x(v), x(v) + |T_v| - 1] \times [y(v) - |T_v| + 1, y(v)]$.



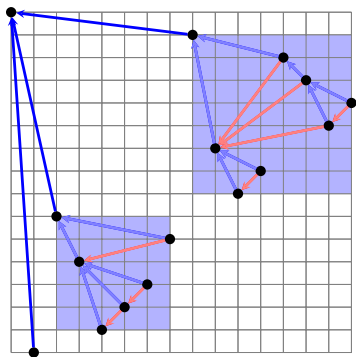
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Edges from v to its children: no crossings!

POLYLINE ROOK-DRAWING FOR PLANAR GRAPHS

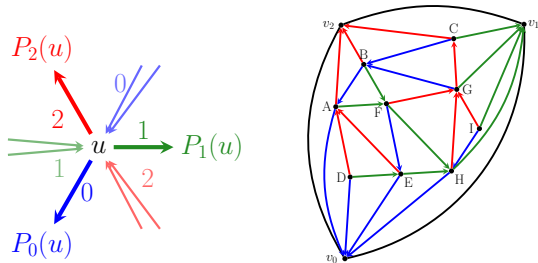
Main result

Every planar graph with n vertices admits a planar polyline rook-drawing, with at most $n - 3$ bends (at most one per edge). Such a drawing can be computed in linear time.

G a triangulation (else, make it triangulated) with exterior vertices v_0 , v_1 and v_2

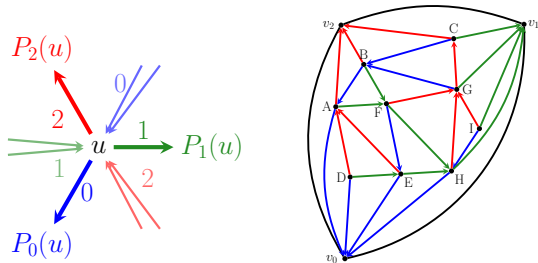
SCHNYDER WOODS

A **Schnyder wood** is a partition of the internal edges of a triangulation in three trees T_0 , T_1 and T_2 (directed toward the root) and with a particular configuration around each inner vertex:



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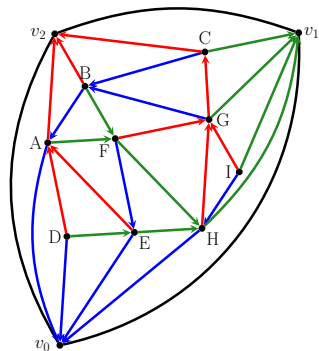


[Schnyder 1989]

Every plane triangulation admits at least one Schnyder wood, and it can be computed in linear time.

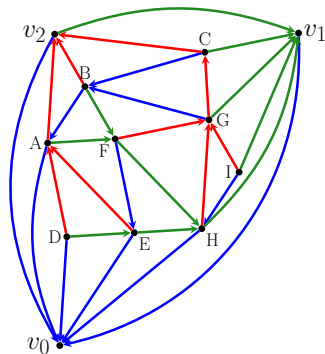
PLANAR POLYLINE ROOK-DRAWING - VERTICES

- (T_0, T_1, T_2) : Schnyder wood of G .



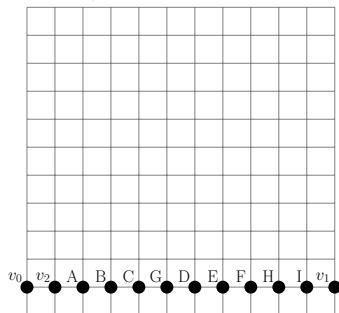
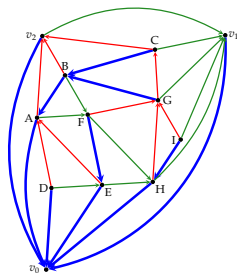
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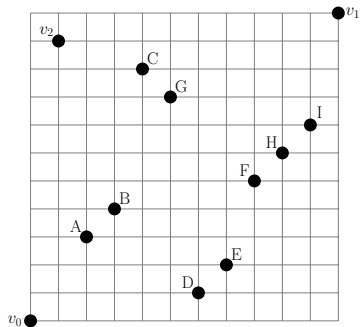
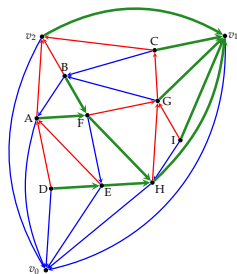
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- $(v_1v_0), (v_2v_0), (v_2v_1)$
- x : clockwise preordering of $T_0 = \{v_0v_2ABC GDEFHIv_1\}$.

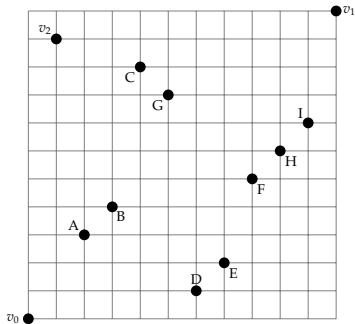
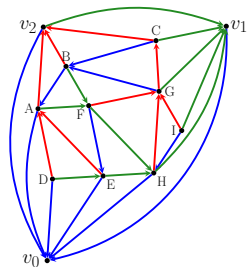


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- y : clockwise postordering of $T_1 = \{DEABFHIGCv_2v_1\}$
($v_0 = 0$).

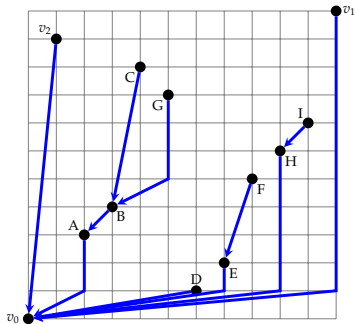
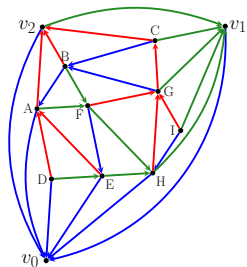


PLANAR POLYLINE ROOK-DRAWING - EDGES



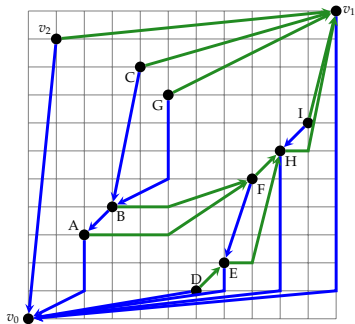
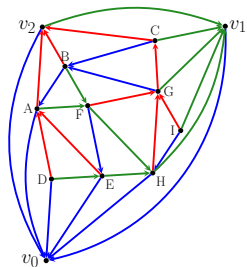
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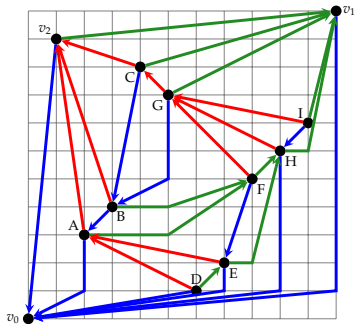
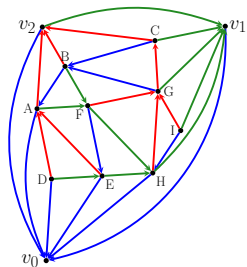
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- Edges of T_2 : not bent



CONCLUSION

Open questions:

- Is a sublinear number of bends sufficient to draw any plane graph planarly?
- If G is a graph with no triangle outer face, what are the conditions to draw G planarly?
- What is the minimum grid size requested to draw a planar straight-lines rook-drawing for a given plane graph? Is this minimum a constant?

Thank you for your attention!